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# Formalisation of Dynamic Properties of Multi-Issue Negotiation

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## **bid\_alternation( $\gamma$ :trace)**

Over time the bids of A and B alternate: thus for all two different moments in time  $t_1, t_3$ , that A generated a bid, there is a moment in time  $t_2$ , with  $t_1 < t_2 < t_3$ , such that A received a bid generated by B.

$\forall A, B: \text{AGENT}, \forall b_1, b_3: \text{BID}, \forall t_1, t_3:$

$t_1 < t_3 \ \&$

$\text{state}(\gamma, t_1, \text{output}(A)) \models \text{to\_be\_communicated\_to\_by}(b_1, B, A) \ \&$

$\text{state}(\gamma, t_3, \text{output}(A)) \models \text{to\_be\_communicated\_to\_by}(b_3, B, A) \Rightarrow$

$\exists b_2, \exists t_2: t_1 < t_2 < t_3 \ \&$

$\text{state}(\gamma, t_2, \text{input}(A)) \models \text{communicated\_to\_by}(b_2, A, B)$

## **is\_followed\_by( $\gamma$ :trace, A:AGENT, t1:time, b1:BID, B:AGENT, t2:time, b2:BID)**

In a negotiation process  $\gamma$  bid  $b_1$  at time  $t_1$  is followed by a bid  $b_2$  at time  $t_2$  iff bids  $b_1$  and  $b_2$  are subsequent bids in  $\gamma$ .

$\text{state}(\gamma, t_1, \text{output}(A)) \models \text{to\_be\_communicated\_to\_by}(b_1, A, B) \ \&$

$\text{state}(\gamma, t_2, \text{output}(B)) \models \text{to\_be\_communicated\_to\_by}(b_2, B, A) \ \&$

$t_1 < t_2 \ \&$

$[ \forall t_3, \forall C, D: \text{AGENT}, \forall b_3: \text{BID}:$

$t_1 < t_3 < t_2 \Rightarrow \text{state}(\gamma, t_3, \text{output}(C)) \not\models \text{to\_be\_communicated\_to\_by}(b_3, C, D) ]$

## **agent\_consecutively\_bids\_to( $\gamma$ :trace, A:AGENT, t1:time, b1:BID, t2:time, b2:BID, B:AGENT)**

In a negotiation process  $\gamma$  agent A consecutively bids  $b_1$  at time  $t_1$  and then  $b_2$  at time  $t_2$  to agent B.

$\text{state}(\gamma, t_1, \text{output}(A)) \models \text{to\_be\_communicated\_to\_by}(b_1, A, B) \ \&$

$\text{state}(\gamma, t_2, \text{output}(A)) \models \text{to\_be\_communicated\_to\_by}(b_2, A, B) \ \&$

$t_1 < t_2 \ \&$

$[ \forall t_3, \forall b_3: \text{BID}:$

$t_1 < t_3 < t_2 \Rightarrow \text{state}(\gamma, t_3, \text{output}(A)) \not\models \text{to\_be\_communicated\_to\_by}(b_3, A, B) ]$

## **stop\_criterion( $\gamma$ :trace, A:AGENT, t2:time)**

The stop criterion holds for agent A at time  $t$ , if at time  $t$  agent A receives a bid by negotiation partner B that is at least as good as the last bid made by A.

$\exists t_1, \exists B: \text{AGENT}, \exists b_1, b_2: \text{BID}:$

$\text{state}(\gamma, t_2, \text{input}(A)) \models \text{communicated\_to\_by}(b_2, A, B) \ \&$

$\text{state}(\gamma, t_1, \text{output}(A)) \models \text{to\_be\_communicated\_to\_by}(b_1, B, A) \ \&$

$\text{is\_followed\_by}(\gamma, t_1, b_1, t_2, b_2) \ \&$

$\text{util}(\gamma, A, b_1) \leq \text{util}(\gamma, A, b_2)$

**negotiation\_continuation( $\gamma$ :trace)**

For both A and B, unless the stop criterion holds, a new proposal is generated by A upon receipt of a proposal by B.

$\forall t, \forall A, B: \text{AGENT}, \forall b1: \text{BID}$ :

$\neg \text{stop\_criterion}(\gamma, A, t) \ \&$

$\text{state}(\gamma, t, \text{input}(A)) \models \text{communicated\_to\_by}(b1, A, B) \Rightarrow$

$[ \exists b2: \text{BID} \exists t2: t2 > t \ \& \ \text{state}(\gamma, t2, \text{output}(A)) \models \text{to\_be\_communicated\_to\_by}(b2, B, A) ]$

**strictly\_dominates( $b1$ :BID,  $b2$ :BID, A:AGENT, B:AGENT)**

A bid  $b1$  dominates a bid  $b2$  with respect to agents A and B iff both agents prefer bid  $b1$  over bid  $b2$ .

$\forall vA1, vA2, vB1, vB2 : \text{real} :$

$\text{util}(A, b1, vA1) \ \& \ \text{util}(A, b2, vA2) \ \& \ \text{util}(B, b1, vB1) \ \& \ \text{util}(B, b2, vB2) \Rightarrow$

$vA1 > vA2 \ \& \ vB1 > vB2$

**weakly\_dominates( $b1$ :BID,  $b2$ :BID, A:AGENT, B:AGENT)**

A bid  $b1$  dominates a bid  $b2$  with respect to agents A and B iff both agents prefer bid  $b1$  over bid  $b2$ .

$\forall vA1, vA2, vB1, vB2 : \text{real} :$

$\text{util}(A, b1, vA1) \ \& \ \text{util}(A, b2, vA2) \ \& \ \text{util}(B, b1, vB1) \ \& \ \text{util}(B, b2, vB2) \Rightarrow$

$vA1 \geq vA2 \ \& \ vB1 \geq vB2$

**strictly\_better\_social\_welfare( $b1$ :BID,  $b2$ :BID, A:AGENT, B:AGENT)**

The social welfare of bid  $b1$  is better than that of bid  $b2$  with respect to agents A and B iff the sum of the utility values of bid  $b1$  is bigger than the sum of the utility values of bid  $b2$ . See also [6,10].

$\forall vA1, vA2, vB1, vB2 : \text{real} :$

$\text{util}(A, b1, vA1) \ \& \ \text{util}(A, b2, vA2) \ \& \ \text{util}(B, b1, vB1) \ \& \ \text{util}(B, b2, vB2) \Rightarrow$

$vA1 + vB1 > vA2 + vB2$

**strictly\_better\_equitability( $b1$ :BID,  $b2$ :BID, A:AGENT, B:AGENT)**

A bid  $b1$  has a better equitability than bid  $b2$  with respect to agents A and B iff the difference in the utility values of bid  $b1$  is less than the difference in utility values of bid  $b2$ .

$\forall vA1, vA2, vB1, vB2 : \text{real} :$

$\text{util}(A, b1, vA1) \ \& \ \text{util}(A, b2, vA2) \ \& \ \text{util}(B, b1, vB1) \ \& \ \text{util}(B, b2, vB2) \Rightarrow$

$| vA1 - vB1 | < | vA2 - vB2 |$

 **$\epsilon$ -equitability( $b$ :BID, A:AGENT, B:AGENT,  $\epsilon$ :real)**

A bid  $b$  has  $\epsilon$ -equitability with respect to agents A and B iff the difference in the utility values of bid  $b$  is less than  $\epsilon$ . Thus, a bid that has an equitability of 0 has a maximum equitability. This definition corresponds to the idea of Raiffa to maximize the minimum utility [10].

$\forall vA, vB : \text{real} :$

$\text{util}(A, b, vA) \ \& \ \text{util}(B, b, vB) \Rightarrow$

$| vA - vB | \leq \epsilon$

**pareto\_inefficiency( $b$ :BID, A:AGENT, B:AGENT,  $\epsilon$ :real)**

With respect to agents A and B, the Pareto inefficiency of a bid  $b$  is the number  $\epsilon$  that indicates the distance to the Pareto Efficient Frontier according to some distance measure  $d$  in utilities. Here  $d(b1, b2)$  is the distance between the bids  $b1$  and  $b2$  when viewed as points in the plane of utilities.

$\forall vA, vB : \text{real} :$

$\text{util}(A, b, vA) \ \& \ \text{util}(B, b, vB) \Rightarrow$

$\text{pareto\_distance}(vA, vB) = \epsilon$

**making\_global\_concession( $\gamma$ :trace, A:AGENT, t1:time, b1:PID, t2:time, b2:PID, B:AGENT)**

In a negotiation process  $\gamma$  agent B makes a global concession to agent B with respect to bid b1 at time t1 and bid b2 at time t2 iff both bids are consecutive, and b2 has a lower utility than b1, from A's perspective. A similar property could be defined stating that an agent receives a global concession from another agent.

agent\_consecutively\_bids\_to( $\gamma$ , A, t1, b1, t2, b2, B) &

$\forall vA1, vA2 : \text{real} :$

util(A, b1, vA1) & util(A, b2, vA2)  $\Rightarrow$

$vA1 > vA2$

**configuration\_differs(b1:PID, b2:PID)**

Two bids b1 and b2 differ in configuration iff there is an issue that has a different value in both bids. Similar properties could be defined stating that two bids differ in configuration in at least x issues.

$\exists a : \text{ISSUE}, \exists v1, v2 : \text{VALUE} :$

value\_of(b1, a, v1) &

value\_of(b2, a, v2) &

$v1 \neq v2$

**agent\_views\_agent\_makes\_config\_variation( $\gamma$ :trace, A:AGENT, B:AGENT, t1:time, b1:PID, t2:time, b2:PID)**

In the view of agent A, agent B varies the configuration, but not the utility. Note that one agent can both be agent A and B, or A and B can refer to different agents.

agent\_consecutively\_bids\_to( $\gamma$ , A, t1, b1, t2, b2, B) &

configuration\_differs(b1, b2) &

$\forall vA1, vA2 : \text{real} :$

util(A, b1, vA1) & util(A, b2, vA2)  $\Rightarrow$

$vA1 = vA2$

**agent\_views\_agent\_makes\_strict\_ε-progression( $\gamma$ :trace, A:AGENT, B:AGENT, t1:time, b1:PID, t2:time, b2:PID, ε:real)**

In the view of agent A, the two consecutive bids b1 and b2 made at times t1 and t2 by agent B show minimum  $\epsilon$ -progression in utility iff the second bid is at least  $\epsilon$  higher than the first bid. Note that one agent can both be agent A and B, or A and B can refer to different agents.

agent\_consecutively\_bids\_to( $\gamma$ , A, t1, b1, t2, b2, B) &

$\forall vA1, vA2 : \text{real} :$

util(A, b1, vA1) & util(A, b2, vA2)  $\Rightarrow$

$vA2 - vA1 > \epsilon$

**strict\_pareto\_monotony( $\gamma$ :trace, tb:time, te:time)**

A negotiation process  $\gamma$  is Strictly Pareto-monotonous for the interval [t1, t2] iff for all subsequent bids b1, b2 in the interval b2 dominates b1:

$\forall t1, t2, \forall A, B : \text{AGENT}, \forall b1, b2 : \text{PID}$

[  $tb \leq t1 < t2 \leq te$  & is\_followed\_by( $\gamma$ , A, t1, b1, B, t2, b2) ]

$\Rightarrow$  strictly\_dominates( $\gamma$ , b2, b1, A, B)

**weak\_pareto\_monotony( $\gamma$ :trace, tb:time, te:time)**

A negotiation process  $\gamma$  is Weakly Pareto-monotonous for the interval [t1, t2] iff for all subsequent bids b1, b2 in the interval b2 weakly dominates b1:

$\forall t1, t2, \forall A, B : \text{AGENT}, \forall b1, b2 : \text{PID}$

[  $tb \leq t1 < t2 \leq te$  & is\_followed\_by( $\gamma$ , A, t1, b1, B, t2, b2) ]

$\Rightarrow$  weakly\_dominates( $\gamma$ , b2, b1, A, B)